

Chapter 2. Connectivity by Charles

Cut Vertex and Cut Edge

Definition

$v \in V(G)$. If $w(G-v) > w(G)$, then
 v is a cut vertex of G .

(w is the # of connected components)

• 该边的顶点不连.

Theorem

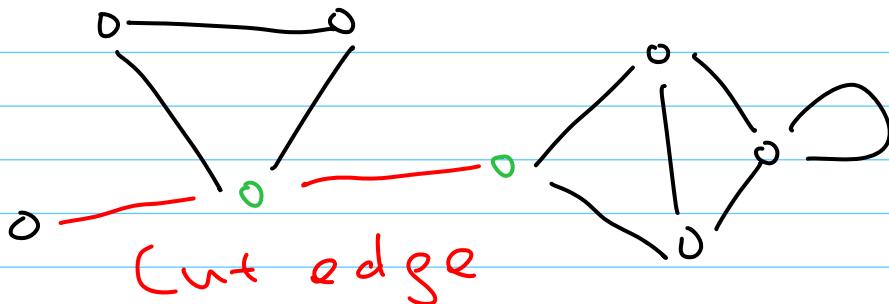
v is cut vertex of tree T if and
only if $d(v) > 1$

Definition

$e \in E(G)$. If $w(G-e) > w(G)$, then
 e is a cut edge of G .

Example:

Cut Vertex



Theorem.

Edge e is cut edge of G if and only if e is contained in no cycle of G

Theorem:

A connected graph is tree if and only if its every edge is cut edge.

Connectivity

Definition:

A vertex cut of G is a subset V' of V such that $G - V'$ is disconnected. A k -vertex cut is a vertex cut of k elements. The connectivity $\kappa(G)$ of G is the minimum k for which G has k -vertex cut.

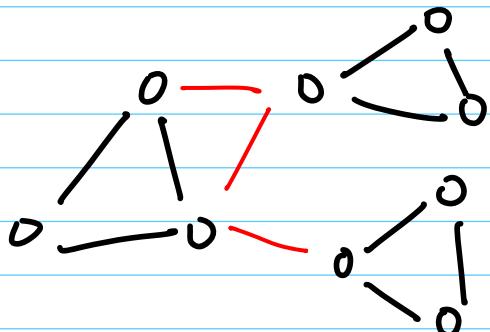
- Cut vertex is 1-vertex cut
- Complete graph has not cut vertex

If $k(G) \geq k$, then G is called
 k -connected

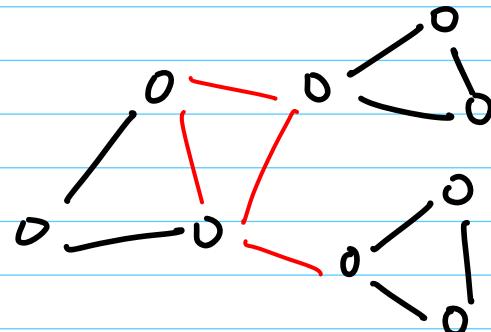
Definition.

An edge cut of G is a subset of E of the form $[S, \bar{S}]$, where S is a nonempty proper subset of V . A k -edge cut a cut edge of k elements. The edge-connectivity of G is the minimum k for which G has k -edge cut: $k'(G)$

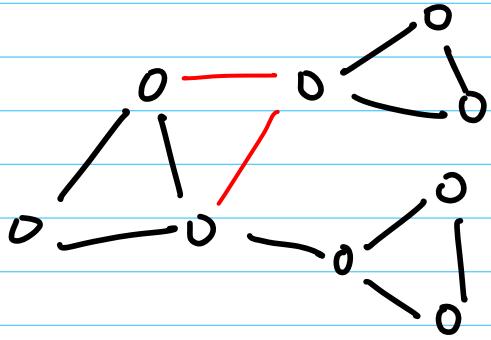
Example



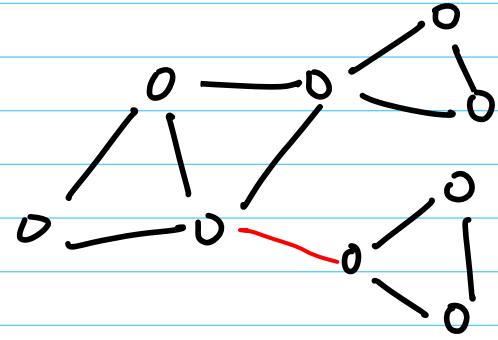
Edge Cut



Not edge Cut



Minimal E-C



Minimum E-C

Theorem

$$k(G) \leq k'(G) \leq s(G)$$

Theorem

Connected Graph with n vertices

and e edges , we have $k(G) \leq \left\lfloor \frac{2e}{n} \right\rfloor$

Corollary

G is a simple Graph, if $s(G) \geq \frac{n-1}{2}$,
then G is connected graph.

Block

Definition:

A connected graph that has no cut vertices is called a block . A

block of a graph is a subgraph that is a block and is maximal with respect to this property.

- Every block with at least three vertices is 2-connected.

Theorem (Whitney, 1932)

If G is 2-connected, then any two vertices of G lie on a common cycle. (are connected by at least two internally-disjoint paths.)

Proof.

← Easy

→ $\exists \text{ 矛法}$

$d(u, v) = 1$. $G - uv$ is connected

because $k' \geq k \geq 2$. \exists Path P

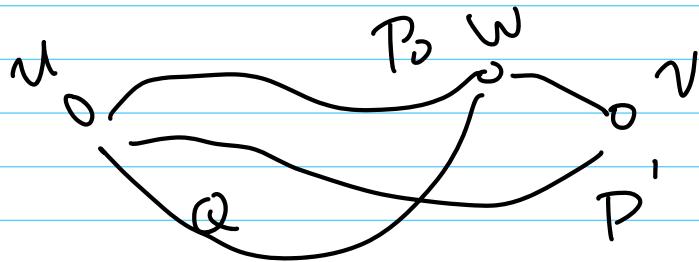
in $G - uv$. $\therefore \exists$ cycle $P + uv$

Assume $d(u, v) < k$ ✓

Let's prove $d(u, v) = k$. ✓

$P_0 = u \dots w v$ with k length

$\therefore d(u, v) = k - 1$; u, v in a cycle,



$\because G - \{w\}$ connected.

$\therefore \exists P'(u, v)$

We can find such a cycle ✓ □

Corollary.

If G is a block with $v \geq 3$,

then any two edges of G lie
on a common cycle.